

Lesson 10:

Relations

- Definitions
- Properties
- Combining Relations
- Relation Closure

R	a	b	c
0	1	1	0
1	1	0	0
2	0	0	0
3	0	1	0

Chapter 7

Relations

Relations are structures that express a connection between 2 or more discrete objects.

- Family Name - Phone number
- Employee - Salary
- Number - Prime factors
- Programs - input parameters

Binary Relations - relate pairs of discrete objects.

Relations

Functions relate between pairs of objects:

$$f : A \rightarrow B$$

Every $a \in A$ is assigned an element $b \in B$.

Cartesian Product relates between pairs of objects:

$$A \times B$$

Every $a \in A$ is related to every $b \in B$.

Binary Relations - relate pairs of elements in A and in B without any constraints.

Relations

Definition:

An *ordered n-tuple* (n-יה סדורה) (a_1, a_2, \dots, a_n) is an ordered collection in which a_1 is first, a_2 is second, ... and a_n as it's n-th element.

2 n-tuples are equal iff they have equal corresponding elements.

$$(a_1, a_2, \dots, a_n) = (b_1, b_2, \dots, b_n) \text{ iff } a_i = b_i \text{ for } i=1, \dots, n$$

An *ordered pair* (זוג סדור) is an n-tuple with $n = 2$.

Relations

Definition:

A *Binary Relation* (יחס (רלציה) בינארית) from set A to set B is a subset of $A \times B$.

$R \subseteq A \times B$ is a relation from A to B then :

$(a,b) \in R$ $a R b$ “a is related to b by R”

“a נמצא ביחס R ל- b”

$(a,b) \notin R$ $a \not R b$ “a is not related to b by R”

Relations

Example:

$$A = \{0, 1, 2\} \quad B = \{a, b\}$$

$$R = \{(0, a), (0, b), (1, a), (2, b)\}$$

R is a relation from A to B because $R \subseteq A \times B$.

$$0 R a \qquad (1, a) \in R$$

$$2 \not R a \qquad (1, b) \notin R$$

Relations

Example:

A = Cities in the world

B = Countries in the World

R is the relation of all (a,b) in which a \in A is the capital of b \in B.

$\left. \begin{array}{l} \{\text{Jerusalem, Israel}\} \\ \{\text{Paris, France}\} \end{array} \right\} \in R$

$\left. \begin{array}{l} \{\text{Jerusalem, USA}\} \\ \{\text{TLV, Japan}\} \end{array} \right\} \notin R$

Relations

Example:

R is a relation on the set of integers.

$$R \subseteq \mathbb{Z} \times \mathbb{Z}$$

$$R = \{ (a,b) \mid a \leq b \}$$

$$(2,100) \in R$$

$$(20,10) \notin R$$

$$(-33, 10,002) \in R$$

$$(0.1, 2) \notin R$$

Representing Relations

$$R \subseteq A \times B.$$

I. List

$$A = \{0, 1, 2, 3\} \quad B = \{a, b, c\}$$

$$R = \{(0, a), (0, b), (1, a), (3, b)\}$$

II. Array of size 2 x 1

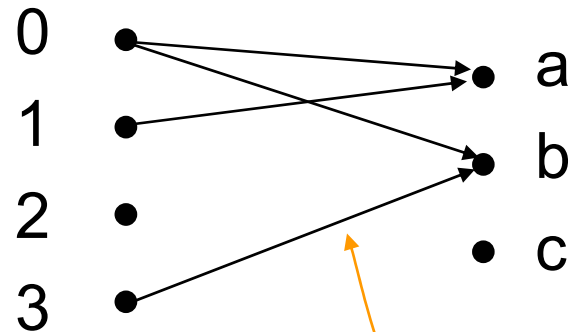
$$A = \{0, 1, 2, 3\} \quad B = \{a, b, c\}$$

$$R = \begin{pmatrix} 0 & 0 & 1 & 3 \\ a & b & a & b \end{pmatrix}$$

Representing Relations

III. Graph

$$A = \{0, 1, 2, 3\} \quad B = \{a, b, c\} \quad R \subseteq A \times B.$$



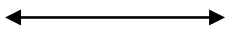
$$R = \{(0, a), (0, b), (1, a), (3, b)\}$$

Representing Relations

III. Table / Binary Matrix

$$A=\{0, 1, 2, 3\} \quad B=\{a, b, c\} \quad R \subseteq A \times B.$$

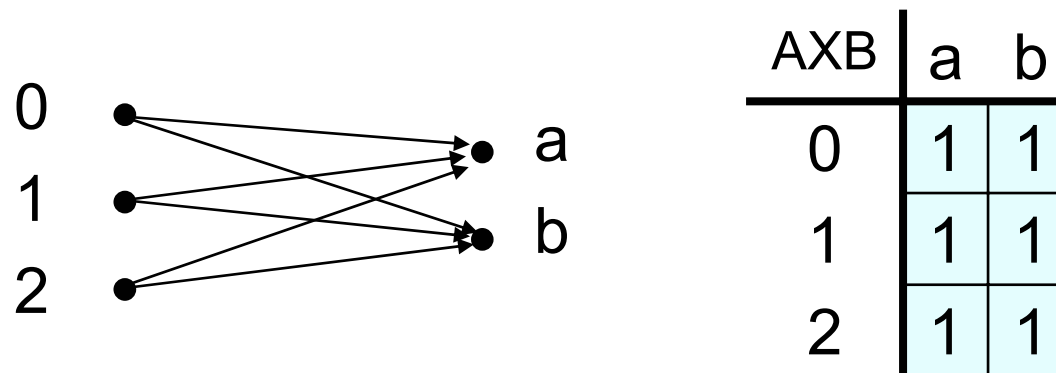
R	a	b	c
0	1	1	0
1	1	0	0
2	0	0	0
3	0	1	0


$$\begin{pmatrix} 1 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix}$$

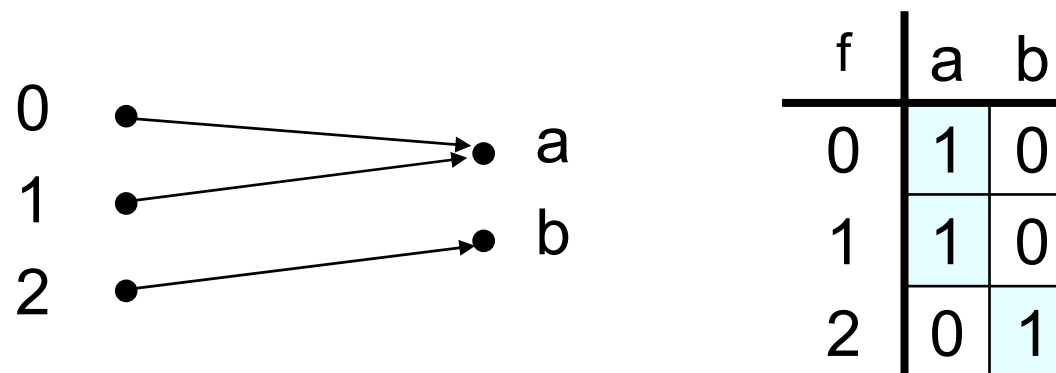

$$R=\{(0, a), (0, b), (1, a), (3, b)\}$$

Relations

The Cartesian Product $A \times B$ is a relation:



Function $f : A \rightarrow B$ is a relation



Relations

Definition:

A relation **on** the set A is a relation from A to A .

Example:

$$A = \{1, 2, 3, 4\}$$

$$R = \{ (a,b) \mid a \mid b \}$$

$$R = \{ (1,1), (1,2), (1,3), (1,4), (2,2), (2,4), (3,3), (4,4) \}$$

$$R \subseteq \mathbb{Z} \times \mathbb{Z}$$

$$R = \{ (a, b) \mid a < b \} \quad \text{infinite set}$$

Relations

Example:

How many relations can be defined on a set A of size n ?

Answer:

A relation is a subset of $A \times A$.

$$|A \times A| = n^2$$

Thus there are 2^{n^2} possible relations.

Relations

Definition:

The *Identity Relation* (יחס היחידה) over the set A ,

$I_A \subseteq A \times A$ is defined as:

$$I_A = \{(a, a) \mid a \in A\}$$

Example:

$$A = \{0, 1\}$$

$$I_A = \{(0, 0), (1, 1)\}$$

Relations

Definition:

The *Inverse Relation* (יחס הפכי) of $R \subseteq A \times B$ is the relation $R^{-1} \subseteq B \times A$ such that:

$$R^{-1} = \{(b, a) \mid (a, b) \in R\}$$

Examples:

$$A = \{0, 1, 2, 3\} \quad B = \{a, b, c\}$$

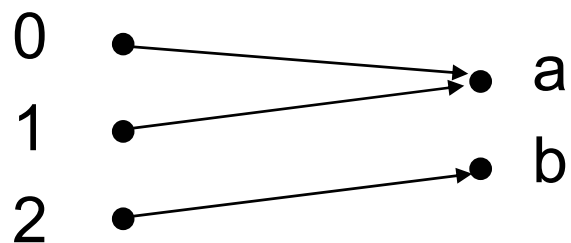
$$R = \{(0, a), (0, b), (1, a), (3, b)\}$$

$$R^{-1} = \{(a, 0), (b, 0), (a, 1), (b, 3)\}$$

Relations

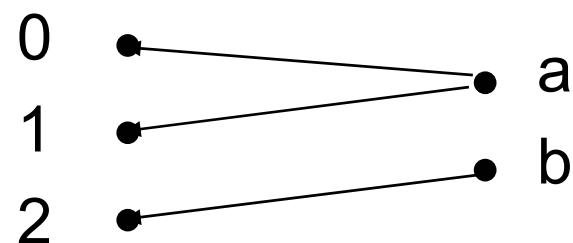
Examples:

$$R = \begin{pmatrix} 1 & 1 & 2 & 3 \\ a & c & c & d \end{pmatrix}$$



$$\begin{matrix} & a & b \\ \begin{matrix} 0 \\ 1 \\ 2 \end{matrix} & \begin{pmatrix} 1 & 0 \\ 1 & 0 \\ 0 & 1 \end{pmatrix} \end{matrix}$$

$$R^{-1} = \begin{pmatrix} a & c & c & d \\ 1 & 1 & 2 & 3 \end{pmatrix}$$



$$\begin{matrix} & 0 & 1 & 2 \\ \begin{matrix} a \\ b \end{matrix} & \begin{pmatrix} 1 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} \end{matrix}$$

Properties of Relations - Reflexivity

Definition:

A relation R on a set A is *reflexive* (רפלקסיבי (חוזר))
iff $\forall a \in A, (a, a) \in R$.

A relation R is reflexive iff $I_A \subseteq R$.

Examples:

$R = \{(x, y) \mid x, y \in \mathbb{Z} \text{ and } x \geq y\}$ reflexive

R is the relation over people s.t. reflexive
 $(x, y) \in R$ if x and y have the same parents.

$R = \{(1, 1), (1, 2), (2, 1)\}$ not reflexive

Properties of Relations - Symmetry

Definition:

A relation R on a set A is *symmetric* (סימטרי)

if $\forall a, b \in A, (a, b) \in R \rightarrow (b, a) \in R$

Definition:

A relation R on a set A is *antisymmetric* (אנטי-סימטרי)

if $\forall a, b \in A, (a, b) \in R \wedge (b, a) \in R \rightarrow a = b$

A relation R is symmetric iff $R = R^{-1}$

A relation R is antisymmetric iff $R \cap R^{-1} \subseteq I_A$

Properties of Relations - Symmetry

Examples:

$$R = \{(1, 1), (1, 2), (2, 1)\}$$

symmetric

$$R = \{(x, y) \mid x, y \in \mathbb{Z} \text{ and } x * y = 6\}$$

symmetric

$$R = \{(1, 1), (1, 2), (2, 3), (4, 2)\}$$

antisymmetric

R is the relation over students s.t.

$(x, y) \in R$ if x and y take the same course.

symmetric

R is the relation over students s.t.

$(x, y) \in R$ if x's grade is higher than y.

antisymmetric

Properties of Relations - Symmetry

The terms symmetric and antisymmetric are not opposites.

$$R = \{(1, 2), (2, 1), (1, 3)\}$$

NOT symmetric
NOT antisymmetric

$$R = \{(1, 1), (2, 2), (3, 3)\}$$

symmetric
& antisymmetric

Properties of Relations - Transitivity

Definition:

A relation R on a set A is *transitive* (טרנזיטיבי (חוזר))
iff $\forall a, b, c \in A, (a, b) \in R \wedge (b, c) \in R \rightarrow (a, c) \in R$.

Examples:

$$R = \{(1, 1), (3, 1), (1, 2)\}$$

not transitive

R is the relation over students s.t.

$(x, y) \in R$ if x 's grade is higher than y .

transitive

$$R = I_A$$

transitive

Properties of Relations

$R \subseteq \mathbb{Z} \times \mathbb{Z}$

	Reflexive	Symmetric	Anti-symmetric	Transitive
$R_1 = \{(x, y) \mid x \leq y\}$	✓	X	✓	✓
$R_2 = \{(x, y) \mid x > y\}$	X	X	✓	✓
$R_3 = \{(x, y) \mid x = y\}$	✓	✓	✓	✓
$R_4 = \{(x, y) \mid x = y \text{ or } x = -y\}$	✓	✓	X	✓
$R_5 = \{(x, y) \mid x = y + 1\}$	X	X	✓	X
$R_6 = \{(x, y) \mid x + y \leq 3\}$	X	✓	X	X
$R_7 = \{(x, y) \mid x \mid y\}$	✓	X	✓	✓
$R_8 = \{(x, y) \mid x+y=6 \wedge x>0\}$	X	X	X	X

Properties of Relations

Reflexive

$$\begin{matrix} & \begin{matrix} a & b & c \end{matrix} \\ \begin{matrix} a \\ b \\ c \end{matrix} & \begin{pmatrix} 1 & & \\ & 1 & \\ & & 1 \end{pmatrix} \end{matrix}$$

M is diagonal

Symmetric

$$\begin{matrix} & \begin{matrix} a & b & c \end{matrix} \\ \begin{matrix} a \\ b \\ c \end{matrix} & \begin{pmatrix} 1 & 1 & \\ 1 & & 1 \\ & 1 & \end{pmatrix} \end{matrix}$$

$M = M^T$

AntiSymmetric

$$\begin{matrix} & \begin{matrix} a & b & c \end{matrix} \\ \begin{matrix} a \\ b \\ c \end{matrix} & \begin{pmatrix} 1 & 1 & \\ 0 & & 0 \\ & 1 & \end{pmatrix} \end{matrix}$$

M is anti-symmetric
elements may or may
not be on the diagonal

Properties of Relations

Question:

How many reflexive relations can be defined on a set A of n elements?

Answer:

There are 2^{n^2} relations that can be defined on A (the number of subsets of $A \times A$).

Reflexive relations contain (a,a) for all $a \in A$.

The rest of the $n^2 - n = n(n-1)$ elements may or may not be in the relation thus the number of reflexive relations on A is:

$$2^{n(n-1)}$$

Properties of Relations

Example: Reflexive relations

$$n = 2 \qquad 2^{2 \cdot 1/2} = 4$$

1	1
1	1

1	1
0	1

1	0
1	1

1	0
0	1

Properties of Relations

Question:

How many symmetric relations can be defined on a set A of n elements?

Answer:

Consider the matrix representation of the relation.

A symmetric matrix is defined by the upper triangle entries + the diagonal entries. (the symmetric entries are then automatically defined).

The number of entries in the upper triangle of the nxn Matrix is:

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

Thus the number of symmetric relations is $2^{\frac{n(n+1)}{2}}$

Properties of Relations

Example: Symmetric relations

$$n = 2 \quad 2^{2 \cdot 3/2} = 8$$

1	1
1	1

0	1
1	1

1	1
1	0

0	1
1	0

1	0
0	1

0	0
0	1

1	0
0	0

0	0
0	0

Properties of Relations

Question:

How many anti-symmetric relations can be defined on a set A of n elements?

Answer:

Consider the matrix representation of the relation.

In an antisymmetric matrix the diagonal may or may not be 1:
 2^n possibilities for the diagonal.

The upper triangle has $1 + 2 + 3 + \dots + (n-1) = \frac{n(n-1)}{2}$ entries.

Each entry 3 possibilities: $M_{ij} = 1 \quad M_{ji} = 0$

$M_{ij} = 0 \quad M_{ji} = 1$

$M_{ij} = 0 \quad M_{ji} = 0$

Thus the number of anti-symmetric relations is $2^n 3^{\frac{n(n-1)}{2}}$

Properties of Relations

Example: Anti-symmetric relations

$$n = 2 \qquad 2^2 * 3^{1*2/2} = 4 * 3 = 12$$

1	1
0	1

0	1
0	1

1	1
0	0

0	1
0	0

1	0
1	1

0	0
1	1

1	0
1	0

0	0
1	0

1	0
0	1

0	0
0	1

1	0
0	0

0	0
0	0

Combining Relations

Theorem: Let R_1 and R_2 be relations from A to B then the following relations are well defined:

$$R_1 \cup R_2, R_1 \cap R_2, R_1 \oplus R_2, R_1 - R_2, R_2 - R_1$$

Combining Relations

Using Matrix representation:

$$M_1 = \begin{pmatrix} 1 & 0 & 1 \\ 1 & 1 & 1 \\ 0 & 0 & 1 \end{pmatrix} \quad M_2 = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 0 & 1 \end{pmatrix}$$

$$R_1 \cup R_2 \leftrightarrow M_1 \vee M_2 = \begin{pmatrix} 1 & 0 & 1 \\ 1 & 1 & 1 \\ 1 & 0 & 1 \end{pmatrix}$$

$$R_1 \cap R_2 \leftrightarrow M_1 \wedge M_2 = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

$$R_1 \oplus R_2 \leftrightarrow M_1 \oplus M_2 = \begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 1 \\ 1 & 0 & 0 \end{pmatrix}$$

Combining Relations

Using Matrix representation:

$$M_1 = \begin{pmatrix} 1 & 0 & 1 \\ 1 & 1 & 1 \\ 0 & 0 & 1 \end{pmatrix} \quad M_2 = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 0 & 1 \end{pmatrix}$$

$$R_1 - R_2 \leftrightarrow M_1 - M_2 = \begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 1 \\ 0 & 0 & 0 \end{pmatrix}$$

$$R_2 - R_1 \leftrightarrow M_2 - M_1 = \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 1 & 0 & 0 \end{pmatrix}$$

$$R_1^{-1} \leftrightarrow 1 - M_1 = \begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 0 \\ 1 & 1 & 0 \end{pmatrix}$$

Relation Composition

Definition:

Let R be a relation from A to B

and S be a relation from B to C

then the *Composition* (הרכבה) of R and S is the relation $S \circ R$ of ordered pairs (a,c)

$$S \circ R = \{(a,c) \mid a \in A, c \in C \text{ and } \exists b (a,b) \in R \wedge (b,c) \in S\}$$

Relation Composition

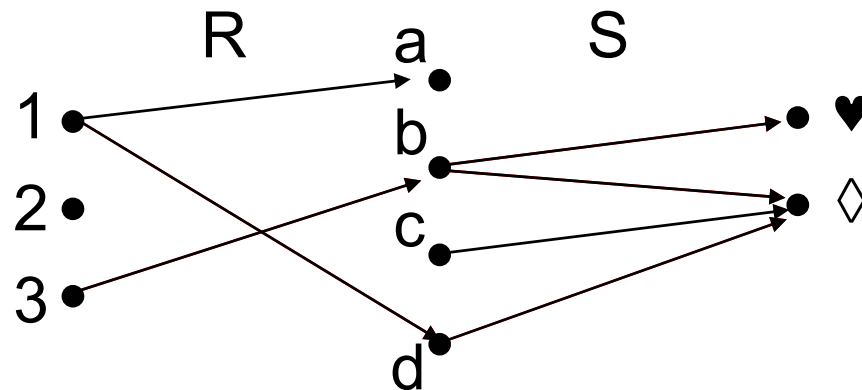
Example:

$$A = \{1,2,3\} \quad B = \{a,b,c,d\} \quad C = \{\heartsuit, \diamond\}$$

$$R \subseteq A \times B \quad R = \{(1,a), (1,d), (3,b)\}$$

$$S \subseteq B \times C \quad S = \{(b,\heartsuit), (b,\diamond), (c,\diamond), (d,\diamond)\}$$

$$S \circ R \subseteq A \times C \quad S \circ R = \{(1,\diamond), (3,\heartsuit), (3,\diamond)\}$$



Relation Composition

Example:

$A = \text{Students in DM}$ $B = \text{Grades (0..100)}$ $C = \text{Pass/nopass}$

$R \subseteq A \times B$ $R = \{ (a,b) \mid b \text{ is the final DM grade of } a \}$

$S \subseteq B \times C$ $S = \{ (b,c) \mid c=\text{Pass for } b < 51 \text{ and}$
 $c=\text{NoPass for } b \geq 51 \}$

$S \circ R \subseteq A \times C$ $S \circ R = \{ (a,c) \mid \text{Student } a \text{ passed the}$
 $\text{DM course } \}$

Relation Composition

$$M_R = \begin{matrix} & a & b & c & d \\ \begin{matrix} 1 \\ 2 \\ 3 \end{matrix} & \begin{pmatrix} \\ \\ \end{pmatrix} \end{matrix}$$

$$M_S = \begin{matrix} & \heartsuit & \diamond \\ \begin{matrix} a \\ b \\ c \\ d \end{matrix} & \begin{pmatrix} \\ \\ \\ \end{pmatrix} \end{matrix}$$

$$M_{S \circ R} = \begin{matrix} & \heartsuit & \diamond \\ \begin{matrix} 1 \\ 2 \\ 3 \end{matrix} & \begin{pmatrix} \\ \\ \end{pmatrix} \end{matrix}$$

$$M_{S \circ R} = M_R \odot M_S$$

Boolean Matrix Product

$$M_{S \circ R}(i,j) = \sum_k M_R(i,k) M_S(k,j) \quad \text{Regular Matrix Product}$$

$$M_{S \circ R}(i,j) = \bigvee_k [M_R(i,k) \wedge M_S(k,j)] \quad \text{Boolean Matrix Product}$$

Relation Composition

$$M_R = \begin{matrix} & a & b & c & d \\ \begin{matrix} 1 \\ 2 \\ 3 \end{matrix} & \begin{pmatrix} \\ \\ \end{pmatrix} \end{matrix}$$

$$M_S = \begin{matrix} & \heartsuit & \diamond \\ \begin{matrix} a \\ b \\ c \\ d \end{matrix} & \begin{pmatrix} \\ \\ \\ \end{pmatrix} \end{matrix}$$

$$M_{S \circ R} = \begin{matrix} & \heartsuit & \diamond \\ \begin{matrix} 1 \\ 2 \\ 3 \end{matrix} & \begin{pmatrix} \\ \\ \end{pmatrix} \end{matrix}$$

$$M_{S \circ R} = M_R \odot M_S$$

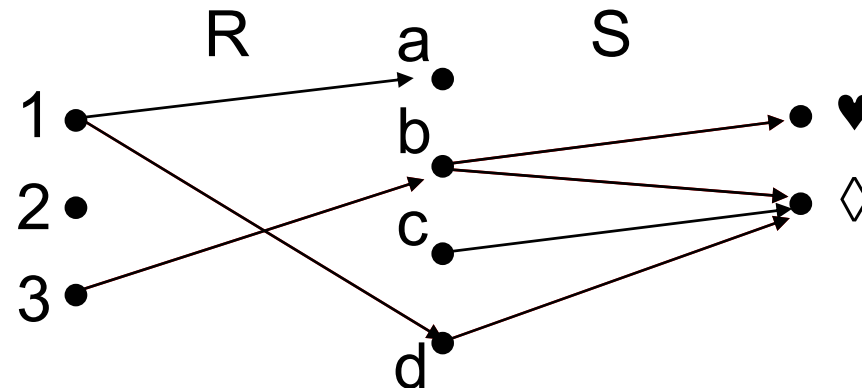
Boolean Matrix Product

$$M_{S \circ R}(a,c) = \sum_b M_R(a,b) * M_S(b,c) \quad \text{Regular Matrix Product}$$

$$M_{S \circ R}(a,c) = \bigvee_b [M_R(a,b) \wedge M_S(b,c)] \quad \text{Boolean Matrix Product}$$

Relation Composition

Example:



$$M_R = \begin{matrix} & \begin{matrix} a & b & c & d \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \end{matrix} & \begin{pmatrix} 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{pmatrix} \end{matrix}$$

$$M_S = \begin{matrix} & \begin{matrix} \heartsuit & \diamond \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \end{matrix} & \begin{pmatrix} 0 & 0 \\ 1 & 1 \\ 0 & 1 \\ 0 & 1 \end{pmatrix} \end{matrix}$$

$$M_{S \circ R} = \begin{pmatrix} 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{pmatrix} \odot \begin{pmatrix} 0 & 0 \\ 1 & 1 \\ 0 & 1 \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 0 & 0 \\ 1 & 1 \end{pmatrix}$$

Relation Composition

Definition:

Let R be a relation on set A .

The **Power n** ($n - \text{תקרת}$) of R ($n \in \mathbb{Z}^+$) is defined recursively:

$$R^1 = R$$

$$R^n = R^{n-1} \circ R$$

$$R^1 = R, R^2 = R \circ R, R^3 = (R \circ R) \circ R \dots$$

Relation Composition

Example:

$$R = \{ (1,1), (2,1), (3,2), (4,3) \}$$

$$R^2 = \{ (1,1), (2,1), (3,1), (4,2) \}$$

$$R^3 = \{ (1,1), (2,1), (3,1), (4,1) \}$$

$$R^4 = \{ (1,1), (2,1), (3,1), (4,1) \}$$

$$R^5 = R^6 = R^7 = \dots \quad \text{for all } n \geq 3$$

Relation Composition

Example:

$$R = \{ (1,1), (2,1), (3,2), (4,3) \}$$

$$M_R = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

$$M_{R^2} = M_R \circ M_R = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{pmatrix}$$

$$M_{R^3} = M_{R^2} \circ M_R = M_R \circ M_{R^2} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{pmatrix}$$

Relation Composition

Example:

$$M_{R4} = M_{R3} \circ M_R = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{pmatrix}$$

$$M_{R5} = M_{R6} = M_{R7} = \dots$$

Combining Relations

Theorem: Relation R on set A is Transitive IFF $R^n \subseteq R$.

Proof:

1) prove $R^n \subseteq R \rightarrow R$ transitive :

$R^n \subseteq R$ so $R^2 \subseteq R$.

If $(a,b) \in R$ and $(b,c) \in R$ then by definition of composition

$(a,c) \in R^2$ however $R^2 \subseteq R$ thus $(a,c) \in R$.

Transitivity is proven.

Combining Relations

Proof cont.:

2) prove R transitive $\rightarrow R^n \subseteq R$: Proof by induction.

$$R^1 = R \subseteq R.$$

Assume $R^n \subseteq R$ prove $R^{n+1} \subseteq R$

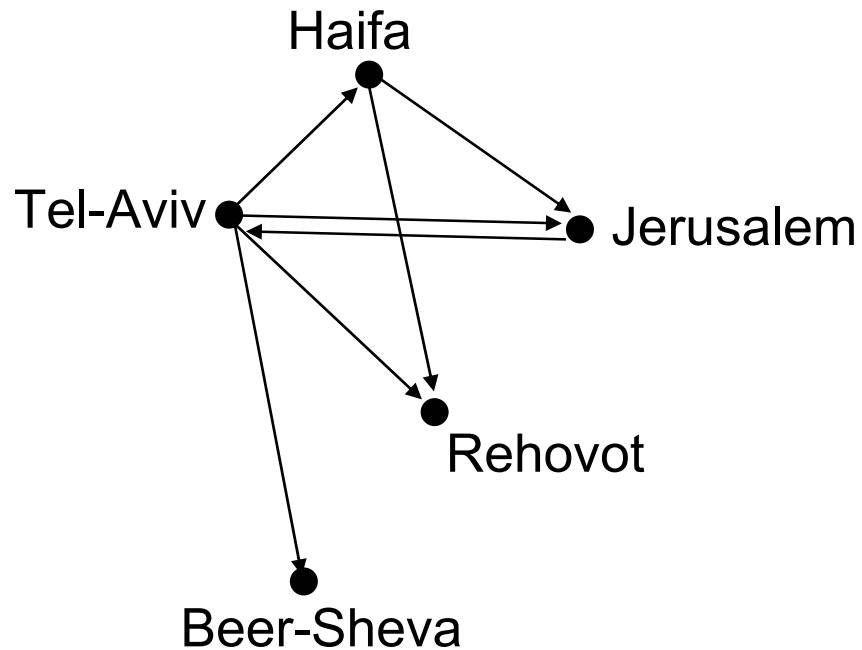
Let $(a,b) \in R^{n+1}$ and show $(a,b) \in R$.

$R^{n+1} = R^n \circ R$ so there exists $x \in A$ such that

$(a,x) \in R$ and $(x,b) \in R^n$, however by the induction assumption $R^n \subseteq R$ and so $(x,b) \in R$.

R is transitive so $(a,x) \in R \wedge (x,b) \in R \rightarrow (a,b) \in R$.

Closure of Relations



Communication Network

R is a relation: $(a, b) \in R$ if a is directionally connected to b .

$(\text{Tel-Aviv}, \text{Jerusalem}) \in R$

$(\text{Haifa}, \text{Tel-Aviv}) \notin R$

Closure of Relations

Definition:

Let R be a relation on set A .

The **Closure** (סגור) of R in terms of property P is the relation S with property P s.t. $R \subseteq S$ and for any other relation $T \neq S$: $R \subseteq T \rightarrow S \subseteq T$.

S is the minimal relation with property P that contains R .

Reflexive Closure

Example:

R is a relation over $A = \{1,2,3\}$

$R = \{ (1,1), (1,2), (2,1), (3,1) \}$

The Reflexive closure of R is obtained by adding:

$(2,2) (3,3)$

Note: any reflexive relation S over A must contain

$(1,1), (2,2), (3,3)$.

Proof ?

Reflexive Closure

The reflexive closure of relation R on A is:

$$R \cup \{(a,a) \mid a \in A\}$$

Example:

$$R = \{ (a,b) \mid a < b \}$$

$$R \subseteq \mathbb{Z} \times \mathbb{Z}$$

$$R = \{ (a,b) \mid a \leq b \}$$

Reflexive closure of R

Symmetric Closure

Example:

R is a relation over $A = \{1,2,3\}$

$R = \{ (1,1), (1,2), (2,3), (3,1), (3,2) \}$

The Symmetric closure of R is obtained by adding:

$(2,1) (1,3)$

Note: any symmetric relation S over A that includes R must contain $(2,1), (1,3)$.

Proof ?

Symmetric Closure

The symmetric closure of relation R on A is:

$$R \cup R^{-1} = R \cup \{(b,a) \mid (a,b) \in R\}$$

Example:

$$R = \{(a,b) \mid a < b\} \quad R \subseteq \mathbb{Z} \times \mathbb{Z}$$

$$R = R \cup \{(a,b) \mid a > b\} \quad \text{Reflexive closure of } R$$

$$= \{(a,b) \mid a \neq b\}$$

Transitive Closure

Example:

R is a relation over $A = \{1,2,3,4,5\}$

$$R = \{ (1,2), (2,3), (3,4), (4,5) \}$$

Add all (a,c) s.t. $(a,b), (b,c) \in R$

$$R \cup \{ (1,3), (2,4), (3,5) \}$$

Transitive?

$$= \{ (1,2), (2,3), (3,4), (4,5), (1,3), (2,4), (3,5) \}$$

Transitive Closure

Example Cont:

$$R' = \{ (1,2), (2,3), (3,4), (4,5) (1,3), (2,4), (3,5) \}$$

Add all (a,c) s.t. $(a,b), (b,c) \in R'$

$$R'' = R' \cup \{ (1,4), (2,5) \}$$

Transitive?

$$R''' = R'' \cup \{ (1,5) \}$$

Transitive!

Transitive Closure

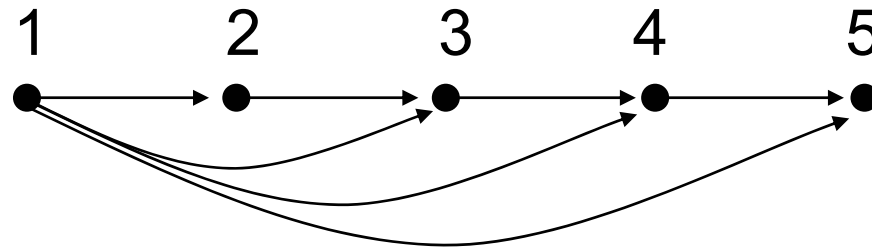
Consider the graph representation of R.

$$R = \{ (1,2), (2,3), (3,4), (4,5) \}$$

Add (1,3)

Add (1,4)

Add (1,5)



Add (1,x) for all x that is 'reachable' from 1.
i.e. there is a 'path' from 1 to 'x'.

Transitive Closure

Definition:

Let R be a relation on set A .

A *Path* (לול) in R from $a \in A$ to $b \in A$ is a sequence of elements of A :

$$x_0 = a, x_1, x_2, \dots, x_n = b$$

$$\text{s.t. } (x_i, x_{i+1}) \in R \text{ for } i = 0 \dots n-1$$

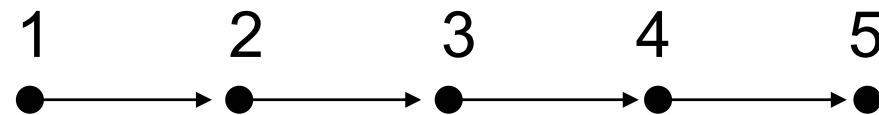
The *length* of the path is n .

Transitive Closure

Example:

R is a relation over $A = \{1,2,3,4,5\}$

$R = \{ (1,2), (2,3), (3,4), (4,5) \}$



There is a path from 1 to 4: 1, 2, 3, 4

There is a path from 2 to 5: 2, 3, 4, 5

There is **no** path from 5 to 3.

Transitive Closure

Theorem: Let R be a relation on set A .

There is a path of length n ($n \in \mathbb{Z}^+$) from a to b IFF
 $(a,b) \in R^n$.

Proof: By induction on n .

Transitive Closure

Transitive closure of R on set A is obtained by adding all (a,b) s.t. there is a path (of any length) from a to b .

Add all pairs (a,b) with a path of length ≤ 2 : R^2

Add all pairs (a,b) with a path of length ≤ 3 : R^3

Add all pairs (a,b) with a path of length ≤ 4 : R^4



Transitive Closure

Theorem: Let R be a relation on set A .
The transitive closure of R is:

$$R^* = \bigcup_{n=1}^{\infty} R^n$$

Proof: Use previous theorem.

Transitive Closure

In practice there is no need to unify ∞ sets R^n .

Theorem: Let R be a relation on set A , $|A|=n$.
If there is a path between a and b in R then
there is a path between a and b of length $\leq n$.

If $a \neq b$ then there is a path of length $n-1$.

Transitive Closure

Proof:

Assume $a, b \in A$ and a path exists between a and b .
Assume, on the contrary, that the shortest path is of length $m > n-1$:

$$a = x_0, x_1, x_2, \dots, x_m = b$$

there are $m+1 > n$ elements of A in the path.

From the Pigeonhole principle 2 of these elements must be equal: $x_i = x_j$. Create a new path:

$$a = x_0, x_1, x_i, x_{j+1}, \dots, x_m = b$$

with at least 1 element less - contradiction that the shortest path is m .

Transitive Closure

In matrix representation: $R = \{ (1,3), (2,1), (3,2) \}$

$$M_R = \begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix}$$

$$R^* = R^1 \cup R^2 \cup R^3$$

$$M_{R^2} = M_R \odot M_R = \begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix} = \begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{pmatrix}$$

$$M_{R^3} = M_{R^2} \odot M_R = \begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{pmatrix} \begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

$$M_{R^*} = M_R \vee M_{R^2} \vee M_{R^3} = \begin{pmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{pmatrix}$$