

# Complexity Theory : Exercise 1

Not for submission

March 10, 2015

1. (Reminder of NP and NP completeness) Consider the languages:

$$Clique = \{(G, k) : G \text{ is a graph that has a clique of size } k\}$$

$$Half - Clique = \{G : G \text{ is a graph that has a clique of size } |V|/2\}$$

(a) Show that  $Half - Clique \in NP$ .

(b) Show that  $Clique \leq_L Half - Clique$ .

2. (Composition of logspace functions).

(a) Let  $M_1, M_2$  be machines that run in space  $O(\log n)$  and compute functions  $f_1$  and  $f_2$  respectively. Consider the function  $f(x) = f_2(f_1(x))$ . Show that this function is computable by a machine that runs in space  $O(\log n)$ .

(b) Show that if  $L_1 \leq_L L_2$  and  $L_2 \leq_L L_3$  then  $L_1 \leq_L L_3$ .

(c) Show that if  $L_1 \leq_L L_2$  and  $L_2 \in L$  then  $L_1 \in L$ .

(d) Show that if  $L_1 \leq_L L_2$  and  $L_2 \in NL$  then  $L_1 \in NL$ .

3. Let  $dPATH = \{G, s, t : G \text{ is a directed graph, } s, t \text{ are vertices and there is a path from } s \text{ to } t\}$ . In class we showed that  $dPATH \in NL$ . Show that  $dPATH$  is NL complete. (Hint: we explained this proof informally in class and the purpose of this exercise is for you to go over it and make sure you understand the details).

4. Show that  $TQBF \in PSPACE$ . (Recall that  $TQBF$  is the set of all true quantified boolean formulae). (Hint: we explained this proof informally in class and the purpose of this exercise is for you to go over it and make sure you understand the details).

5. ( $2 - SAT$  is NL complete) Consider the language  $2 - SAT = \{\phi : \phi \text{ is a satisfiable 2-CNF formula}\}$ . (Reminder: a 2-CNF formula is a conjunction of clauses where each clause is the disjunction of two literals).

In class we showed that  $2 - SAT \in NL$ . (Recall that we actually showed that  $2 - SAT \in co - NL$  and used the fact that  $NL = co - NL$ ). Our goal is to show that  $2 - SAT$  is NL complete.

(a) Show that  $co - dPATH = \{G, s, t : G \text{ is a directed graph, } s, t \text{ are vertices and there is no path from } s \text{ to } t\}$  is NL complete.

(b) Show that  $co - dPATH \leq_L 2 - SAT$ . (Hint: use the ideas we used in class to relate  $2 - SAT$  to connectivity in graphs).

(c) Conclude that  $2 - SAT$  is NL complete.